

# DUKE MATH MEET 2012

## TIEBREAKER ROUND SOLUTIONS

1. An 8-inch by 11-inch sheet of paper is laid flat so that the top and bottom edges are 8 inches long. The paper is then folded so that the top left corner touches the right edge. What is the minimum possible length of the fold?

*Solution.* Label the vertices of the rectangle clockwise from upper-left as  $A, B, C, D$ , so that  $A$  gets folded to  $X \in CD$ . Then the fold is the perpendicular bisector  $\ell$  of  $AX$ . Then we have three cases: either  $\ell$  intersects  $AB$  and  $CD$ ,  $\ell$  intersects  $AB$  and  $AD$ , or  $\ell$  intersects  $AD$  and  $BC$ .

In the first case let  $\ell$  intersect  $AB$  at  $P$  and  $CD$  at  $Q$ . Write  $\angle BAX = \theta$ . Then construct  $P' \in AB$  such that  $P'D \parallel PQ$ . Then we have  $P'D = PQ$ . Furthermore, triangle  $ADP'$  is similar to triangle  $BAX$ . Hence we have  $PQ = P'D = AD \sec \theta$ . This is minimized when  $\theta = 0^\circ$ , or when  $A$  is folded to  $B$  and  $PQ = 11$ .

In the second case, let  $\ell$  intersect  $AB$  at  $P$  and  $AD$  at  $Q$ . Let  $PQ$  intersect  $AX$  at  $R$ , which is the midpoint of  $AX$ . Let  $H$  be the foot of the perpendicular from  $B$  to  $AX$ . Then we know that triangle  $AQP$  is similar to triangle  $BAX$ . Hence  $PQ/AR = AX/BH$ . But as  $\angle ABX = 90^\circ$ , we may calculate the area of triangle  $ABX$  in two different ways to get  $AX \cdot BH = AB \cdot BX$ . Hence we have

$$PQ = \frac{AR \cdot AX}{BH} = \frac{AX^2}{2BH} = \frac{AX^3}{2AB \cdot BX}.$$

Writing  $\theta = \angle BAX$ , we find that  $PQ = AB/(2 \sin \theta \cos^2 \theta)$ ; hence we need only maximize  $\sin \theta \cos^2 \theta$ . Writing  $u = \sin \theta$ , we have  $\sin \theta \cos^2 \theta = \sin \theta - \sin^3 \theta = u - u^3$ . Making the substitution  $u = 2 \sin \phi / \sqrt{3}$ , we have  $u - u^3 = 2 \sin(3\phi) / 3\sqrt{3} \leq 2/3\sqrt{3}$ . Hence we find that  $PQ \geq 3\sqrt{3}AB/4 = 6\sqrt{3}$ . (This minimization may also be done more quickly with calculus, but this particular non-calculus-based technique is rather nice.)

In the third case, let  $\ell$  intersect  $BC$  at  $P$  and  $AD$  at  $Q$ . Construct  $Q' \in AD$  such that  $BQ' \parallel PQ$ . Then we know that  $BQ' = PQ$ . We know also that triangle  $AQ'B$  is similar to triangle  $BAX$ , so that  $BQ'/AB = AX/BX$ . Hence to minimize  $BQ' = PQ$ , we need only minimize the ratio  $AX/BX$ . Writing  $\theta = \angle BAX$ , we have  $AX/BX = \csc \theta$ . To minimize  $\csc \theta$ , we want  $\angle BAX$  as large as possible, so we take  $X = C$ . Then this gives  $AX = \sqrt{8^2 + 11^2} = \sqrt{185}$  and  $BX = 11$ . Thus we get  $PQ \geq 8\sqrt{185}/11$ .

In order to determine which case gives us the global minimum, we need to determine the relative ordering of  $8\sqrt{185}/11$ ,  $6\sqrt{3}$ , and  $11$ . As it turns out, we have  $8\sqrt{185}/11 < 6\sqrt{3} < 11$ , so the minimum possible length of the crease is  $8\sqrt{185}/11$ .

2. Triangle  $ABC$  is equilateral, with  $AB = 6$ . There are points  $D, E$  on segment  $AB$  (in the order  $A, D, E, B$ ), points  $F, G$  on segment  $BC$  (in the order  $B, F, G, C$ ), and points  $H, I$  on segment  $CA$  (in the order  $C, H, I, A$ ) such that  $DE = FG = HI = 2$ . Considering all such configurations of  $D, E, F, G, H, I$ , let  $A_1$  be the maximum possible area of (possibly degenerate) hexagon  $DEFGHI$  and let  $A_2$  be the minimum possible area. Find  $A_1 - A_2$ .

*Solution.* We know that  $[DEFGHI] = [ABC] - ([ADI] + [BFE] + [CHG])$ . Hence maximizing  $[ADI] + [BFE] + [CHG]$  is equivalent to minimizing  $[DEFGHI]$ . Write  $u = AD, v = BF, w = CH$ . Then we know that

$$[ADI] + [BFE] + [CHG] = \frac{\sqrt{3}}{4} [u(4-w) + v(4-u) + w(4-v)],$$

where we have  $0 \leq u, v, w \leq 4$ . Clearly the minimum occurs when  $u = v = w = 0$ , so that  $[DEFGHI] = [ABC] = 9\sqrt{3}$ .

For convenience write  $f(u, v, w) = u(4-w) + v(4-u) + w(4-v)$ . Now we claim that  $f(u, v, w) \leq 16$ . We show that for any  $0 \leq u, v, w \leq 4$ , either  $f(0, v, w) \geq f(u, v, w)$  or  $f(4, v, w) \geq f(u, v, w)$ . Indeed we have

$$\begin{aligned} f(0, v, w) - f(u, v, w) &= u(v+w-4); \\ f(4, v, w) - f(u, v, w) &= (u-4)(v+w-4). \end{aligned}$$

As  $u$  and  $u-4$  have opposite signs it follows that one of the two differences will be non-negative. Hence in maximizing  $f$  we may assume that  $u, v, w \in \{0, 4\}$ . To obtain a maximum clearly we cannot have  $u = v = w = 0$  or  $u = v = w = 4$ . But if one or two of  $u, v, w$  are 4 and the others are 0, then  $f(u, v, w) = 16$ . Hence  $f(u, v, w) \leq 16$  for all  $0 \leq u, v, w \leq 4$ .

Thus  $[DEFGHI]$  achieves its minimum when  $[ADI] + [BFE] + [CHG] = 4\sqrt{3}$ , so that  $[DEFGHI] = 9\sqrt{3} - 4\sqrt{3} = 5\sqrt{3}$ . We find  $A_1 - A_2 = 4\sqrt{3}$ .

3. Find

$$\tan \frac{\pi}{7} \tan \frac{2\pi}{7} \tan \frac{3\pi}{7}.$$

*Solution.* By De Moivre's formula, we know that for  $\theta \in \mathbb{R}$  we have

$$(\cos \theta + i \sin \theta)^k = \cos(n\theta) + i \sin(n\theta).$$

Take  $k = 7, \theta = n\pi/7$ , and consider the imaginary parts of both sides. The imaginary part of the right-hand side is zero, while we can find the imaginary part of the left-hand side by the binomial theorem. This gives

$$7 \cos^6 \theta \sin \theta - 35 \cos^4 \theta \sin^3 \theta + 21 \cos^2 \theta \sin^5 \theta - \sin^7 \theta = 0.$$

Dividing by  $\sin^7 \theta$ , which is nonzero, gives

$$7 \cot^6 \theta - 35 \cot^4 \theta + 21 \cot^2 \theta - 1 = 0,$$

which holds for  $\theta = \pi/7, \theta = 2\pi/7$ , and  $\theta = 3\pi/7$ . Thus we have by Vieta's formulas that

$$\cot^2 \frac{\pi}{7} \cot^2 \frac{2\pi}{7} \cot^2 \frac{3\pi}{7} = \frac{1}{7},$$

so inverting and taking square roots gives

$$\tan \frac{\pi}{7} \tan \frac{2\pi}{7} \tan \frac{3\pi}{7} = \sqrt{7}.$$